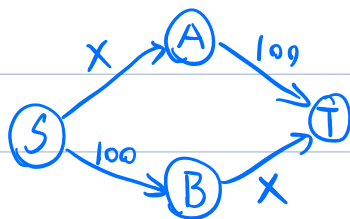


## Today: Algorithmic Game Theory (AGT)

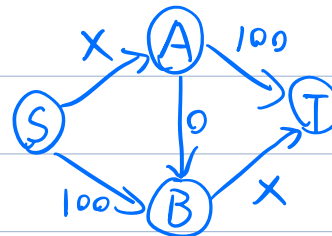
- Game theory x CS
- Algorithms in "strategic" environments

### 1. Traffic Networks

Imagine the following two traffic networks with costs.



Network 1

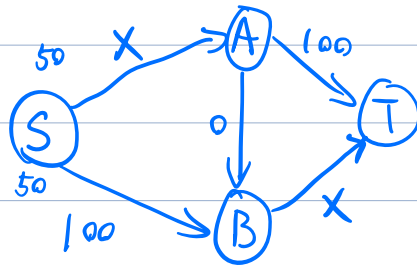


Network 2

For cost  $X$ , it means if there are  $X$  cars going from  $S$  to  $T$ , the cost is equal to  $X$ . Each driver minimizes their own cost, and we want to minimize average cost. Which network is better?

- If single driver, network 2.
- If 100 drivers, and we control where they go, network 2.
- If 100 drivers w/ free will what happens?

With this  $A \rightarrow B$  shortcut, the first driver to take that path  $S \rightarrow A \rightarrow B \rightarrow T$  only has cost  $50 + 50 = 100 < 150$ . More drivers will take this route, when # of drivers reach 50, the cost of new path becomes



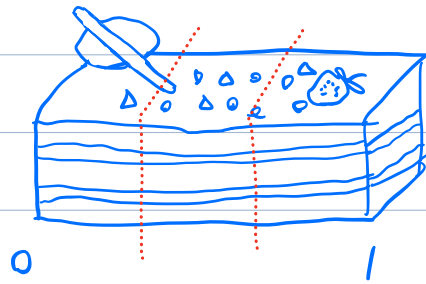
150, same as before, so no more improvements as opposed to S-A-T. But drivers in S-B-T has time 200, so they'll switch to S-A-B-T.

In the end, everyone goes S-A-B-T, average cost is 200! But without A-B, the average was only 150.

### ↑ Braess' Paradox

Key takeaway: when we want to say a network (algorithm/system) is "better", we need to explicitly model how people will use it.

## 2. Cake Cutting / Fair Division



Cutting a cake  $[0, 1]$  and divide among  $n$  players.

Each player has a different evaluation of the cake  $v_i : S \subseteq [0, 1] \rightarrow [0, 1]$ .

$$\text{e.g. } v_i([0, \frac{1}{3}]) = 1$$

$$v_i([\frac{1}{3}, 1]) = 0$$

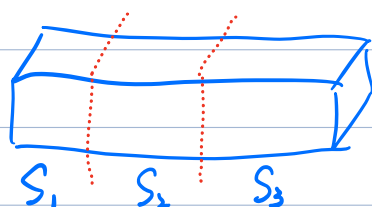
Envy-free: An allocation  $S_1, S_2, \dots, S_n$  to  $n$  players is envy-free if

$$\forall i, j, v_i(S_i) \geq v_i(S_j)$$

Two players: cut and choose

Player A cuts into 2 pieces, player B chooses first

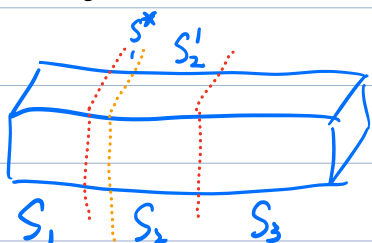
Three players?



Player A cuts into 3 pieces, let B & C choose?

What if they want the same piece? How to ensure fairness among B, C?

Idea: Have B "trim" its favorite piece, say  $S_2$   
to have the same value as its second favorite.



Now there are two  
"favorite pieces" for B,

Now let C choose between  $S_1, S_2', S_3$ .

Then B choose from the rest, and A gets the remaining.

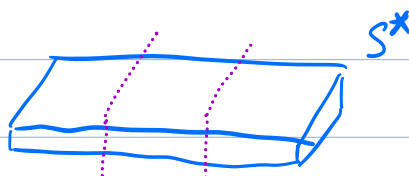
A: 😊

B: 😊

C: 😊

Done if we throw the trimming  $S^*$  away.

what if we want to keep it?



Let's say B gets  $S'_2$ , now we let C cut.

(If C gets  $S'_2$ , we let B cut.)

B chooses first, then A, then C.

B & C both happy, why?

Before allocating the trimming, both are happy.

& one gets to cut, and one gets to choose first.

A also happy.

A's cut in step 1 already better than entire  $S_2$ .

A's cut in step 1 as good as C's cut.

--- step 2 ---



↑ Selfridge and Conway (~1960)

More players? 4 players, 20 cuts, by Aziz and Mackenzie (2016)

$n$  players,  $n^{n^{n^n}}$  cuts, by Aziz and Mackenzie (2016)