

Name:  
Net ID:

**Basic Algorithms (Section 7)**  
Fall 2024

HW2 (Due 9/25 22:00)  
Instructor: Jiaxin Guan

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Problem 1A (Recurrence Relations — 3 Styles, 9 pts)

Use induction to show the asymptotic run-time corresponding to the following recurrence relation is  $O(\log^2(n))$ . **Do not use the Master Theorem here.**

$$T(0) = T(1) = 1. \text{ For all } n \geq 2, T(n) = T(\lceil \frac{n}{2} \rceil - 1) + \log n.$$

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Problem 1B (9 pts)

Use a recursion tree to find the asymptotic run-time corresponding to the following recurrence relation. You may assume any fractional input to  $T$  is the greatest integer less than it (e.g.,  $T(\frac{3n}{5}) = T(\lfloor \frac{3n}{5} \rfloor)$ ).

$$T(0) = T(1) = 1. \text{ For all } n \geq 2, T(n) = T(\frac{3n}{5}) + T(\frac{4n}{5}) + n^2$$

Problem 1C (12 pts)

Find the asymptotic run-time corresponding to each of the following recurrence relations using the Master Theorem. For each, explain which case of the Master Theorem applies and why.

(a)  $T(n) = 3T(\frac{n}{8}) + \sqrt{2n}$ .

(b)  $T(n) = 3T(\frac{n}{3}) + 3n$ .

(c)  $T(n) = 4T(\frac{n}{2}) + n \log^2 n$ .

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Problem 2 (Array Search, 30 pts)

You are given an array of  $n$  integers  $a_1 < a_2 < \dots < a_n$ . Give an  $O(\log n)$  algorithm that outputs an index  $i$  where  $a_i = i$ , or outputs  $\perp$  if such  $i$  does not exist. Justify the correctness and time complexity of your proposed algorithm.

Problem 3 (Malfunctioning Phones, 40 pts)

A manufacturer has a recall on a set of  $n$  cell phones, some of which have a malfunction which makes them unreliable. The manufacturer has built a machine that allows a pair of phones to test each other's correctness. Let  $C_1, C_2$  be a pair of phones. The machine  $M$  runs in the following way:

1.  $M(C_1, C_2) = 11$  if both phones say the other is working.
2.  $M(C_1, C_2) = 10$  if one phone says the other is working and one phone says the other is malfunctioning.
3.  $M(C_1, C_2) = 00$  if both phones say the other is malfunctioning.

Remember that malfunctioning phones cannot be trusted, so they may lie, tell the truth, or throw out a random response. Working phones, on the other hand, can be assumed to know if the other phone is working or malfunctioning always.

- (a) Show that if you know at least one working phone, all other working phones can be found by using  $O(n)$  queries to  $M$ .
- (b) Assume the majority of the phones are working, i.e., there are greater than  $n/2$  working phones. Give an algorithm that can find a working phone in  $O(n)$  queries to  $M$ . Justify the correctness and time complexity of your proposed algorithm. [Hint: Start by explaining how to use  $O(n)$  queries to reduce the problem size by a constant factor.]
- (c) Assume the majority of the phones are malfunctioning, i.e., there are fewer than  $n/2$  working phones. Is there still a procedure (using  $M$ ) that is guaranteed to find a working phone? Give a brief justification.